

Binary Trees and Heaps

CS 491 – Competitive Programming

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Objectives

- ▶ Write code to implement binary search trees, regular trees, and heaps.
- ▶ Explain the differences between these data structures.

Binary Search Trees

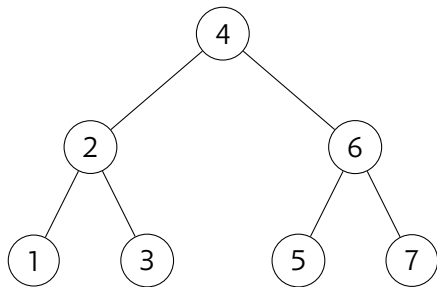
A *Binary Search Tree* is a set of vertices v_i such that:

- ▶ There is precisely one vertex with no parents, called the *root*.
- ▶ Each vertex has 0, 1, or 2 children.
- ▶ Each vertex has a value of a type that supports ordering
- ▶ If a child value is less than the parent, it must be the left child.
- ▶ If a child value is greater than the parent, it must be the right child.
- ▶ If the values are equal: go left, go right, or delete. Just be consistent.

Important notes:

- ▶ Expected height of the tree is $\mathcal{O}(\log_2 n)$.
- ▶ Worst case height is $\mathcal{O}(n)$. When does this occur?

Picture



Implementing

- ▶ Use a struct for simplicity

```
1 struct bst<T> {  
2     T value;  
3     bst<T> *left, *right;  
4  
5     bst<T>(T value) {  
6         this->value = value;  
7         left = right = NULL;  
8     }  
9 }
```

- ▶ You could also be clever and use a sized-2 vector for the children, or a pair.

Add

```
1  bst<T> add(bst<T> *root, T value) {
2      if (root == NULL) {
3          return bst<T>(value);
4      }
5      bst<T> *parent = root;
6      while (true) {
7          if (value < root->value) {
8              if (parent->left == NULL) {
9                  return parent->left = bst<T>(value);
10             } else {
11                 parent = parent->left;
12             }
13         } else {
14             // Same thing, but go right.
15         }
16     }
17 }
```

Find

```
1  bool find(bst<T> *root, T value) {
2      while (root) {
3          if (root->value == value) {
4              return true;
5          }
6          if (value < root->value)
7              root = root->left;
8          else
9              root = root->right;
10     }
11     return false;
12 }
```

Deletion

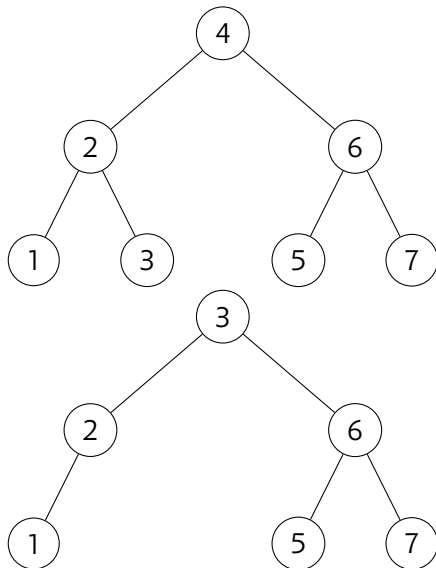
Deletion is a bit of a pain. The steps:

- ▶ Find the victim node
- ▶ Get the In Order Predecessor (IOP) of the victim node.
- ▶ Replace the victim value with the IOP value.
- ▶ Delete the IOP from the child branch.

There are edge cases!

- ▶ Deleting the last vertex
- ▶ Nodes without an IOP

Picture: Deleting 4



Regular Trees

If we don't constrain order or number of children, we just have a tree.
Trees have special properties!

- ▶ $|E| = |V| - 1$. Adding even one more edge makes it not a tree anymore.
- ▶ There are no cycles.
- ▶ Equivalently: there is exactly one path between any two nodes.

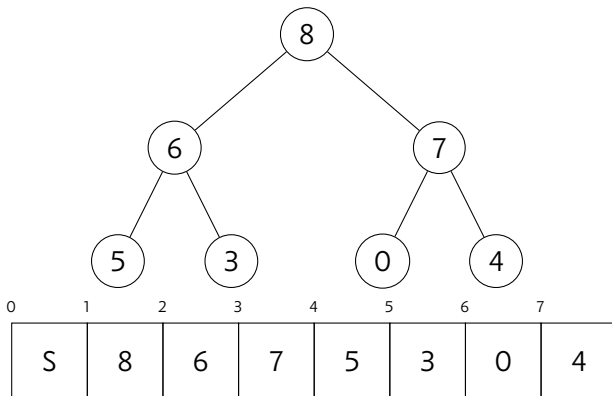
Representation:

- ▶ Usually we would use the adjacency list representation from last time to construct the tree.

Remember Heaps?

- ▶ A heap is also a binary tree.
 - ▶ Each child is smaller (max-heap) or larger (min-heap) than the parent.
- ▶ We use vectors to represent them.
 - ▶ Leave the 0 element empty as a sentinel. The math is cleaner this way.
- ▶ You rarely need to use heaps as heaps, but this method of storing a binary tree is often very efficient!

Heap Visualization



Implementing

```
1  int goLeft(int i) {
2      return 2*i;
3  }
4
5  int goRight(int i) {
6      return 2*i + 1;
7  }
8
9  int goUp(int i) {
10     if (i>1) {
11         return i / 2;
12     }
13 }
```